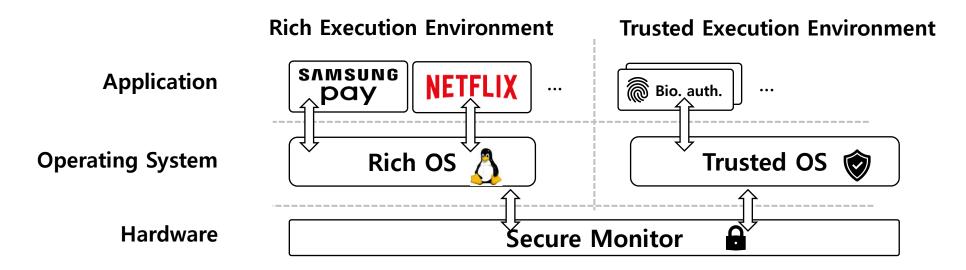
Formal Specification of Trusted Execution Environment APIs

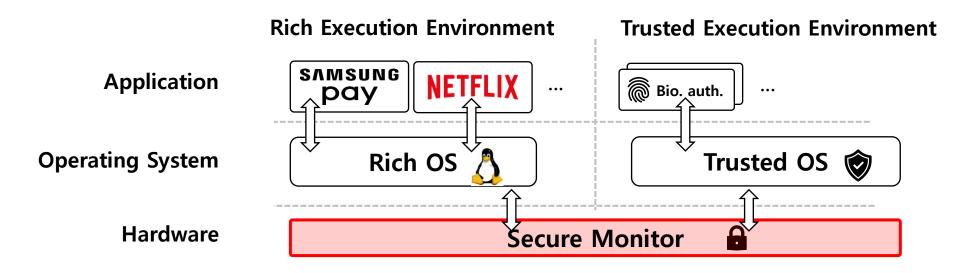
<u>Geunyeol Yu</u>¹ Seunghyun Chae¹ Kyungmin Bae¹ Sungkun Moon²

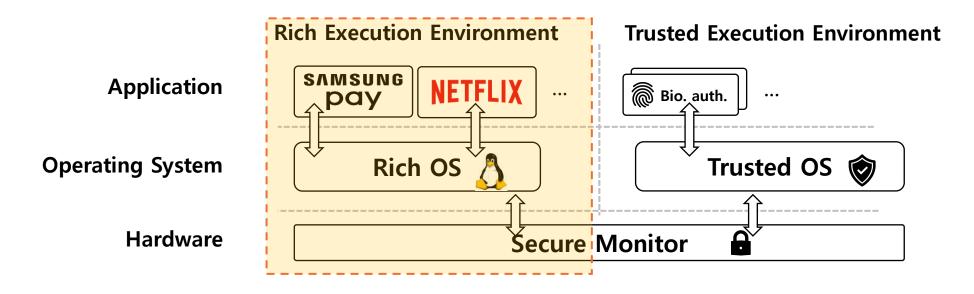
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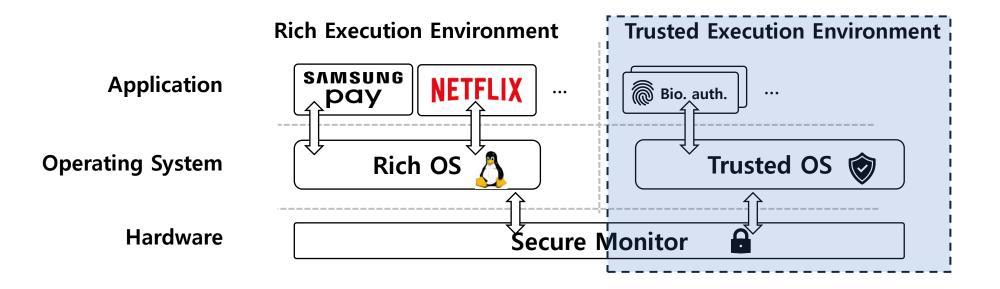


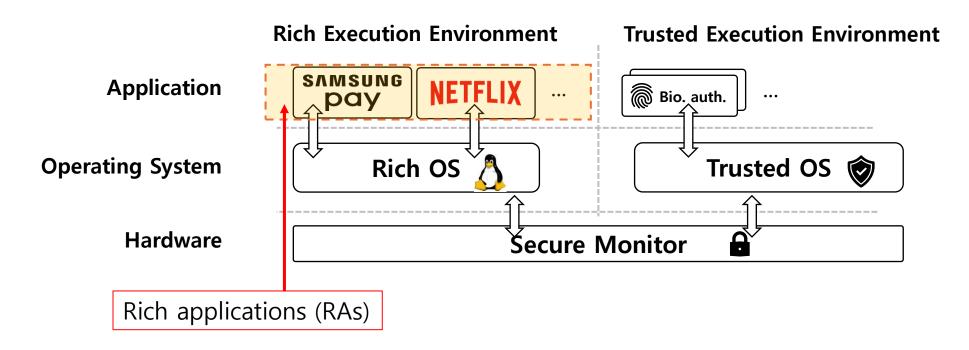


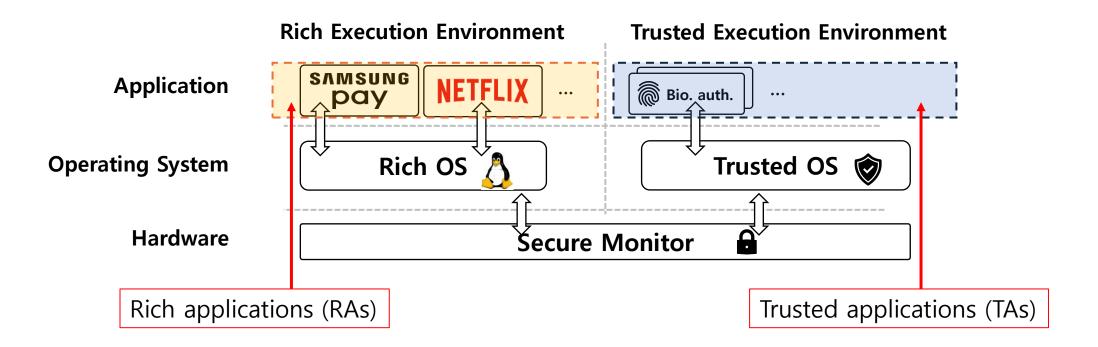


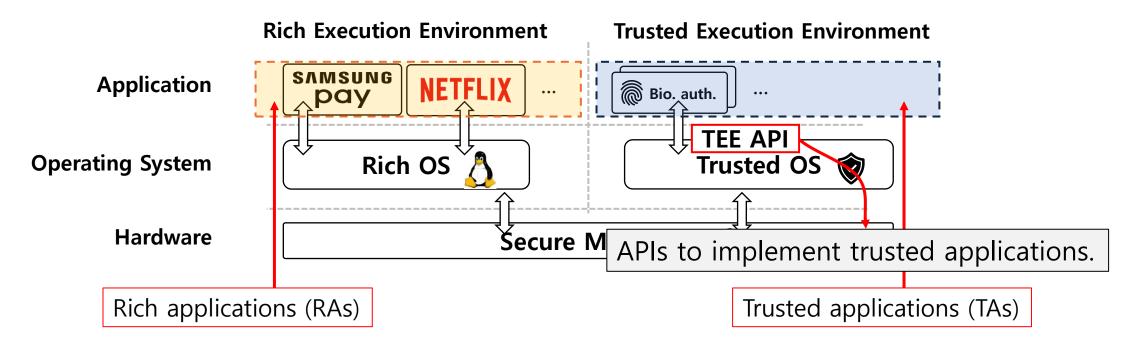












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• This is why TEE is widely used in security-critical systems, such as industrial control systems, servers, mobile security, IoT, etc.

Motivations

• Formal analysis framework for TEE applications is not well-developed.

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• Formal analysis framework for TEE applications is not well-developed.

• Formal models for TEE and its APIs, which can be utilized for a variety of formal analysis techniques, are lacking.

- We provide a comprehensive formal model for <u>TEE APIs</u>, that can be used in various formal analysis.
- We specify two widely used TEE API categories, Trusted Storage API and Cryptographic Operations API.
- We demonstrate the effectiveness of our model through a case study on formally analyzing a real-world TEE application, MQT-TZ.
 - Identify security vulnerabilities in the MQT-TZ implementation.
 - Patch them and verify the fix with model checking.

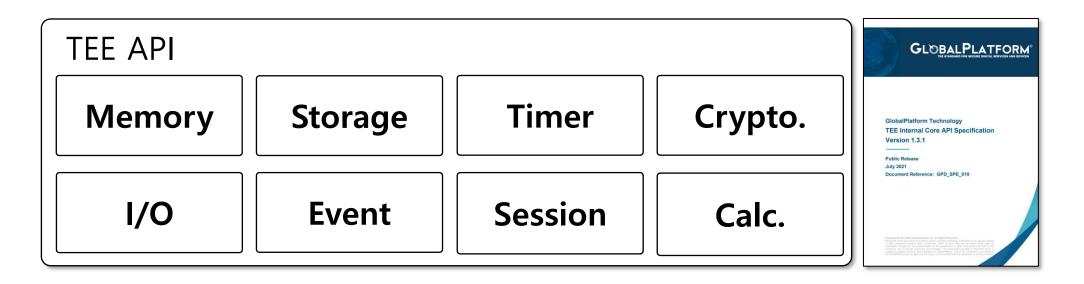
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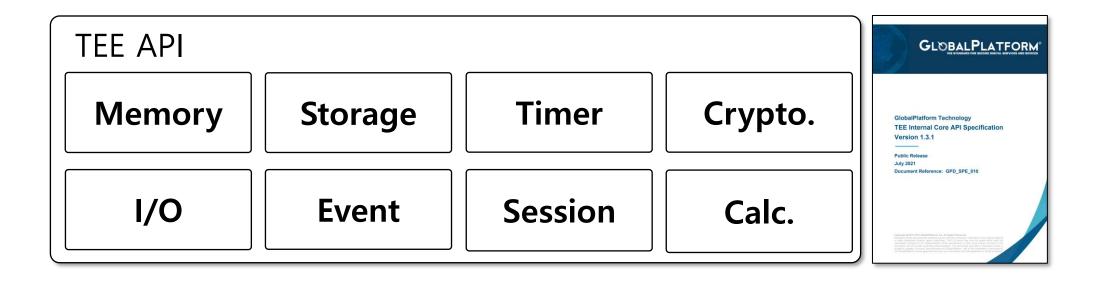
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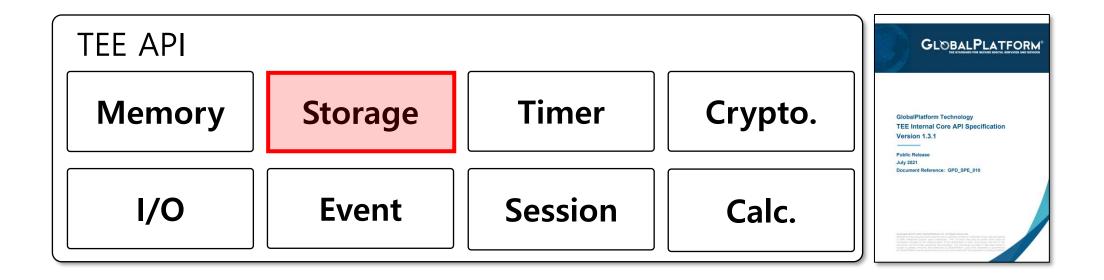
- Our target is the standard TEE APIs, provided by Global Platform.
 - Many Trusted OSes follow this standard.
 - e.g., Samsung TEEgris, Trustonic Kinibi, Qualcomm QTEE, etc.



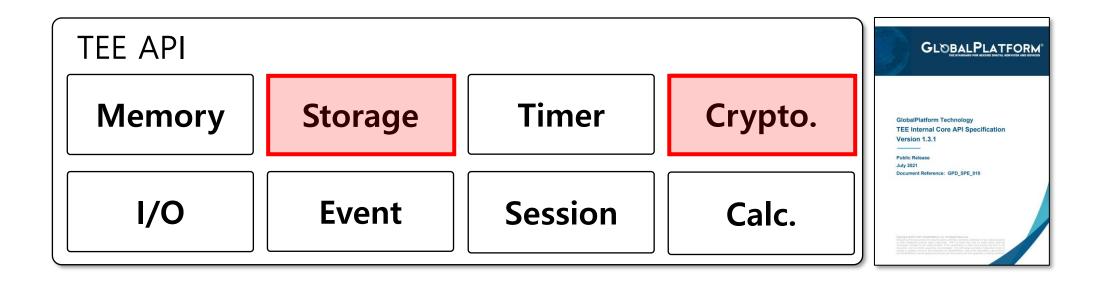
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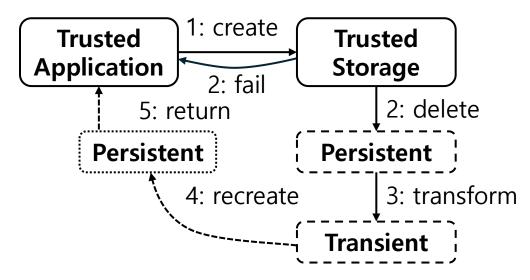
Handles cryptographic algorithms

- We choose these APIs because:
 - They are widely and frequently used in various TEE applications;
 - They provide essential functions for TEE's integrity.

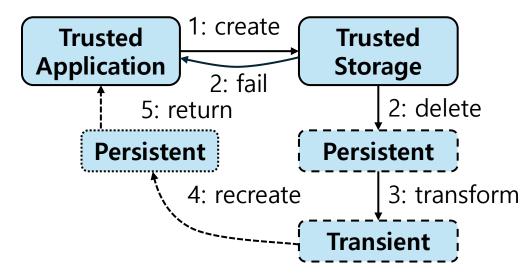
		I/O	Event	Session	Calc.		
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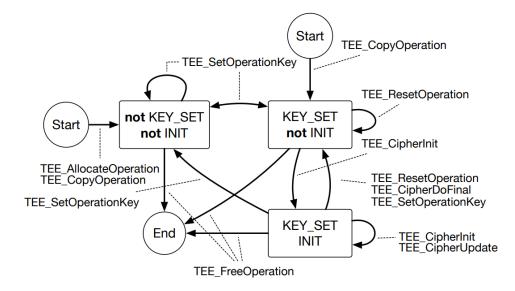
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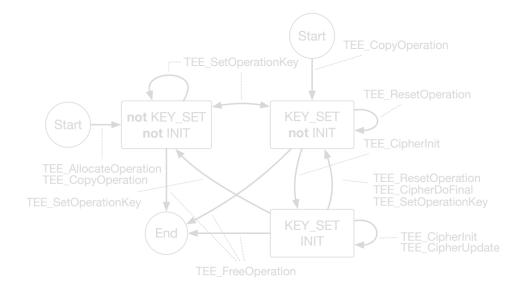
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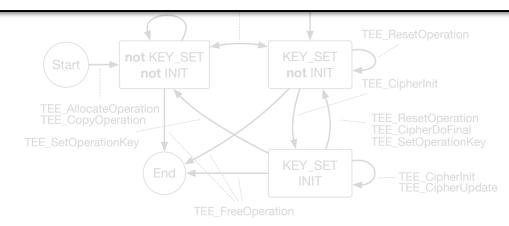
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Considering these characteristics, we use Maude for formal specification.



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We can formally specify TEE APIs considering characteristic 1 and 2.

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 - It supports object-oriented specification.
 - It defines concurrent behaviors using rewrite rules.
- Because of the powerful formalism of Maude, it is widely used in various formal analysis domains such as:
 - defining language semantics,
 - inductive theorem proving,
 - model checking, etc.

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```
rl [read]:
    (msg reqRead from TA to PI)
    < PI : PersistObj | file-name : FILE, data-stream : DATA :: STREAM >
=> < PI : PersistObj | file-name : FILE, data-stream : STREAM >
    (msg retData[DATA] from PI to TK)
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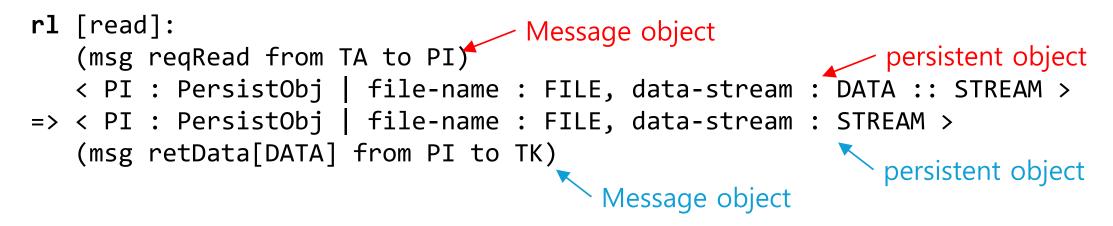
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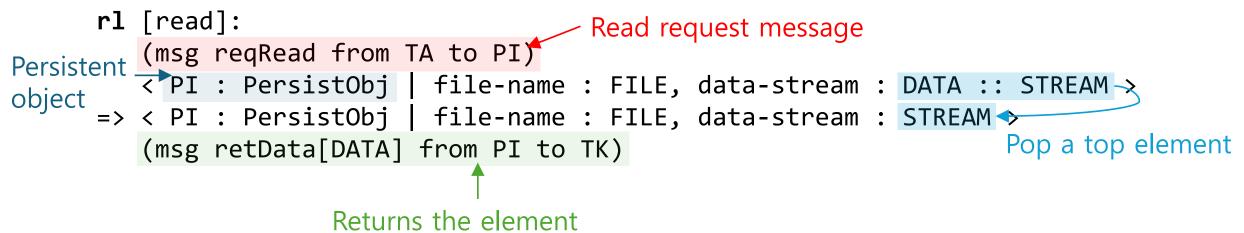
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```
rl [read]:
Persistent
object
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(msg retData[DATA] from PI to TK)
Pop a top element</pre>
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File

• This function creates a new persistent object.

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 - Argument 1 : Filename
 - Argument 2 : Access flags (e.g., overwrite)
 - Argument 3 : Data

- ...

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 - Argument 1 : Filename
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It's a file open function but opens the file to a trusted storage.

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• According to the TEE API document, when a file with the same name already exists, the behavior of the function is as follows:

5.7.2 TEE_CreatePersistentObject Since: TEE Internal Core API v1.3 - See Backward Compatibility note below. TEE_Result TEE_CreatePersistentObject(

TEE Internal Core API Specification - Public Release v1.3.1

uint32_t storageID, [in(objectIDLength)] void* objectID, size_t objectIDLen, uint32 t flags, TEE_ObjectHandle attributes, initialData, size_t initialDataLen, [outopt] TEE_ObjectHandle* object);

Description

The TEE CreatePersistentObject function creates a persistent object with initial attributes and an initial data stream content. The storageID parameter indicates which Trusted Storage Space to access; possible values are defined in Table 5-2.

The flags parameter is a set of flags that controls the access rights, sharing permissions, and object creation mechanism with which the object handle is opened. The value of the flags parameter is constructed by a bitwise-inclusive OR of flags from the following list:

- Access control flags:
- TEE DATA FLAG ACCESS READ: The object is opened with the read access right. This allows the Trusted Application to call the function TEE ReadObjectData.
- · TEE DATA FLAG ACCESS WRITE: The object is opened with the write access right. This allows the Trusted Application to call the functions TEE_WriteObjectData and TEE TruncateObjectData
- TEE_DATA_FLAG_ACCESS_WRITE_META: The object is opened with the write-meta access right. This allows the Trusted Application to call the functions
- TEE CloseAndDeletePersistentObject1 and TEE RenamePersistentObject
- Sharing permission control flags:
- TEE DATA FLAG SHARE READ: The caller allows another handle on the object to be created with read access.
- TEE_DATA_FLAG_SHARE_WRITE: The caller allows another handle on the object to be created with write access
- TEE_DATA_FLAG_OVERWRITE: As summarized in Table 5-13:
- If this flag is present and the object exists, then the object is deleted and re-created as an atomic operation: that is, the TA sees either the old object or the new one
- o If the flag is absent and the object exists, then the function SHALL return TEE ERROR ACCESS CONFLICT.
- Other flags are reserved for future use and SHALL be set to 0

The attributes of the newly created persistent object are taken from attributes, which can be another persistent object or an initialized transient object. The object type, size, and usage are copied from attributes.

To create a pure data object, the attributes argument can also be NULL. If attributes is NULL, the object type SHALL be set to TEE_TYPE_DATA to create a pure data object.

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159/375

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TEE Internal Core API Specification - Public Release v1.3.1 159/375 5.7.2 TEE_CreatePersistentObject Since: TEE Internal Core API v1.3 - See Backward Compatibility note below. TEE Result TEE CreatePersistentObject(uint32_t storageID, [in(objectIDLength)] void* objectID, size_t objectIDLen, uint32 t flags, TEE_ObjectHandle attributes, initialData, size_t initialDataLen, [outopt] TEE_ObjectHandle* object);

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Delete the old file and create a new one

	Internal Core API v1.3		patibility note below.
TEE_	Result TEE_CreateP	ersistentObject(uint32 t	storageID.
54	in(objectIDLength)]		objectID, size t objectIDLen,
2.	in(objectiozengen/j	uint32 t	flags,
		TEE_ObjectHandle	
	[nbuf]	void*	initialData, size_t initialDataLen,
[0	outopt]	TEE_ObjectHandle*	object);
Decembration	-		
Descriptio			
data stream			a persistent object with initial attributes and an in es which Trusted Storage Space to access; poss
The flogs			cess rights, sharing permissions, and object creat value of the flags parameter is constructed b
mechanism	with which the object h usive OR of flags from t		value of the Frags parameter is constructed b
mechanism bitwise-inclu			value of the frags parameter is consuduced b
mechanism bitwise-inclu • Acce	usive OR of flags from t ss control flags:	he following list:	opened with the read access right. This allows t

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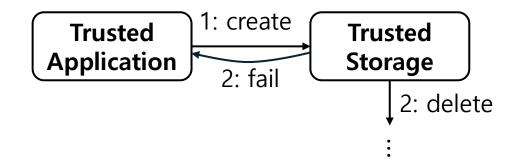
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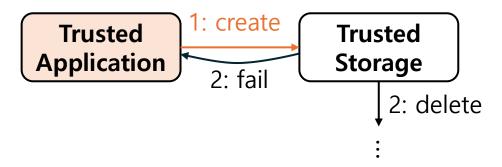
Delete the old file and create a new one

- Overwrite flag not given :

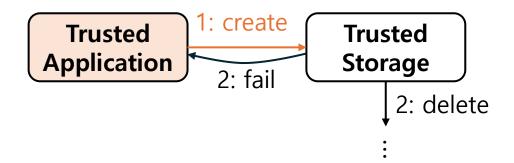
Return error

Since	: TEE Internal Core API v1.3	- See Backward Con	patibility note below.	
	TEE_Result TEE_CreateP			
	[in(objectIDLength)]	uint32_t void* uint32_t TEE ObjectHandle	<pre>storageID, objectID, size_t objectIDLer flags, attributes.</pre>	1,
	[inbuf] [outopt]	void* TEE_ObjectHandle*	<pre>initialData, size_t initialD object);</pre>)ataLen,
Desc	ription			
data s			a persistent object with initial attribu es which Trusted Storage Space to	
mech		handle is opened. The	ccess rights, sharing permissions, an value of the flags parameter is	
•	Access control flags:			
	 TEE_DATA_FLAG_ACCES Trusted Application to car 		opened with the read access right. adObjectData.	This allows
		o call the functions TE	is opened with the write access righ E_WriteObjectData and	t. This allow:
	This allows the Trusted A	Application to call the f	object is opened with the write-meta unctions and TEE_RenamePersistentObje	-
•	Sharing permission control f	lags:		
	 TEE_DATA_FLAG_SHARE read access. 	_READ: The caller allo	ws another handle on the object to	be created w
	 TEE_DATA_FLAG_SHARE with write access. 	_WRITE: The caller al	lows another handle on the object to	be created
•	TEE_DATA_FLAG_OVERWRI	TE: As summarized in	Table 5-13:	
-	 If this flag is present and operation: that is, the TA 		the object is deleted and re-created ject or the new one.	l as an atomi
	 If the flag is absent and t TEE_ERROR_ACCESS_CO 		he function SHALL return	
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persis			e taken from attributes, which e object type, size, and usage a	
	eate a pure data object, the t type SHALL be set to TEE_		nt can also be NULL. If attribut a pure data object.	es is NULL,



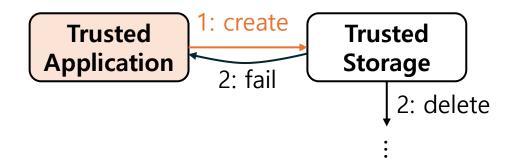


• (1) A trusted application (TA) requests a trusted storage to create a file.



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5.7.2	TEE_CreatePersist	entObject	
Since:	TEE Internal Core API v1.5	- See Backward Con	spatibility note below.
	TEE_Result TEE_CreateP	ersistentObject(wint32 t	storageID.
	[in(object20Length)]		objectID, size_t objectIDLen, flags,
	[inbuf] [outopt]	void* TEE_ObjectHandle*	<pre>initialDwta, size_t initialDwtalen, object);</pre>
Descr	iption		
data st			a persistent object with initial attributes and an initia les which Trusted Storage Space to access; possibil
mecha		handle is opened. The	coss rights, sharing permissions, and object creation value of the flags parameter is constructed by
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	TEE_DATA_FLAG_ACCES Trusted Application to ca		s opened with the read access right. This allows the radObjectData.
		o call the functions TE	is opened with the write access right. This allows $E_W^{-}iteObjectData$ and
	This allows the Trusted A	application to call the f	object is opened with the write-meta access right. unctions and TEE_RenamePersistentObject.
• •	Sharing permission control t	lags:	
	TEE_DATA_FLAG_SHARE read access.	_READ: The caller alk	ows another handle on the object to be created with
	TEE_DATA_FLAG_SHARE with write access.	_NRITE: The caller al	lows another handle on the object to be created
• 1	TEE_DATA_FLAG_OVERNRI	TE: As summarized in	Table 5-13:
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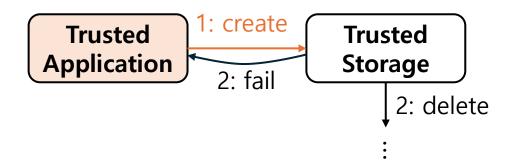


- (1) A trusted application (TA) requests a trusted storage to create a file.
- Trusted application has the following things:
 - the status of an API call,

- ...

- an identifier of a trusted storage,

5.7.2	TEE_CreatePersist	entObject		
Since	TEE Internal Core API v1.3	- See Backward Com	patibility note below.	
	TEE_Result TEE_CreateP	ersistentObject(wint32 t	storageID.	
	[in(objectIDLength)]	void* vint32_t TEE ObjectHandle	objectID, size_t objectIDLen, flags, attributes.	
	[inbuf] [outopt]	void* TEE_ObjectHandle*	initialData, size_t initialDa	italen,
Desci	ription			
data s			a persistent object with initial attribut es which Trusted Storage Space to a	
mecha		handle is opened. The	cess rights, sharing permissions, and value of the flags parameter is o	
•	Access control flags:			
	 TEE_DATA_FLAG_ACCES Trusted Application to ca 		opened with the read access right." adobjectData.	This allows the
		o call the functions TE	is opened with the write access right E_WriteObjectData and	This allows
	This allows the Trusted A	Application to call the fi	object is opened with the write-meta anctions and TEE_RenamePersistentObje	-
	Sharing permission control I	lags:		
	 TEE_DATA_FLAG_SHARE read access. 	_READ: The caller allo	ws another handle on the object to b	e created with
	 TEE_DATA_FLAG_SHARE with write access. 	_WRITE: The caller al	iows another handle on the object to	be created
٠	TEE_DATA_FLAG_OVERNRI	FE: As summarized in	Table 5-13:	
	 If this flag is present and operation: that is, the TA 		the object is deleted and re-created ject or the new one.	as an atomic
	 If the flag is absent and I TEE_ERROR_ACCESS_CO 		he function SHALL return	
•	Other flags are reserved for	future use and SHALL	be set to 0.	
persist			e taken from attributes, which e object type, size, and usage a	
	ate a pure data object, the type SHALL be set to TEE_		nt can also be NULL. If attribute a pure dala object	s is NULL, th
			rreed. revisions, and estensions by GlobalPlathers and our use inconsistent with that servemen	Use of this t is strictly

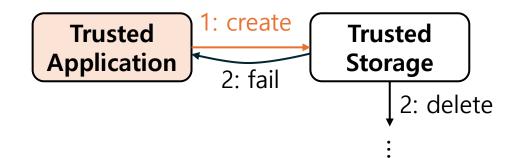


- (1) A trusted application (TA) requests a trusted storage to create a file.
- Trusted application has the following things:
 - the status of an API call,

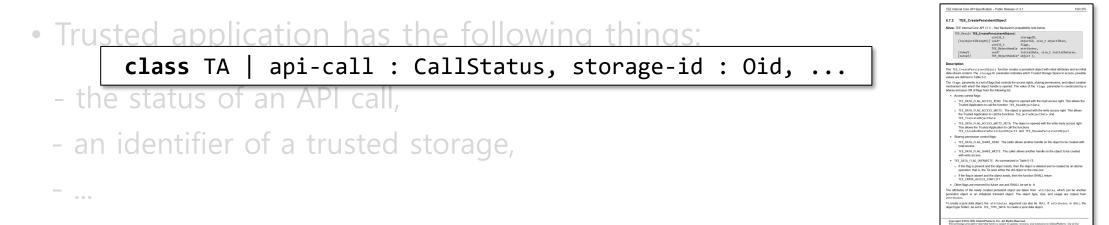
- ...

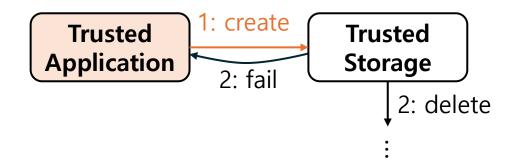
- an identifier of a trusted storage,

5.7.2	TEE_CreatePersist	entObject		
Since: Ti	EE Internal Core API v1.3	- See Backward Corr	patibility note below.	
TE	E_Result TEE_CreateP			
	[in(objectIDLength)]	vint32_t void* vint32_t TEE ObjectMandle	storageID, objectID, size_t objectIDLen, flags, attributes.	
	[inbuf] [outopt]	void* TEE_ObjectHandle*	<pre>initialDuta, size_t initialDutalen, object);</pre>	
Descrip	tion			
data strea			a persistent object with initial attributes and a es which Trusted Storage Space to access; p	
mechanis		andle is opened. The	cess rights, sharing permissions, and object o value of the flags parameter is construct	
 Ap 	cess control flags:			
	TEE_DATA_FLAG_ACCES Trusted Application to ca		opened with the read access right. This allow adObjectData.	vs the
		call the functions TE	is opened with the write access right. This all E_WriteObjectData and	285
	This allows the Trusted A	oplication to call the fi	object is opened with the write-meta access r anctions and TEE_RenamePersistentObject.	ight.
• Sh	aring permission control f	lags:		
	TEE_DATA_FLAG_SHARE read access.	_READ: The caller allo	ws another handle on the object to be create	d with
	TEE_DATA_FLAG_SHARE with write access.	_NRITE: The caller al	iows another handle on the object to be creat	ed
 TE 	E_DATA_FLAG_OVERWRIT	E: As summarized in	Table 5-13:	
	If this flag is present and operation: that is, the TA		the object is deleted and re-created as an ab ject or the new one.	omic
	If the flag is absent and fl TEE_ERROR_ACCESS_CO		he function SHALL return	
 Of 	ner flags are reserved for	future use and SHALL	be set to 0.	
	t object or an initialized		e taken from attributes, which can be a e object type, size, and usage are copie	
	a pure data object, the e SHALL be set to TEE_		nt can also be NULL. If attributes is NU a pure data object.	LL, th
The feeh		ervin is subject to updates.	rred. resisions, and estensions by GlobalPlatform. Use of th and any use incandident with that agreement is strictly	

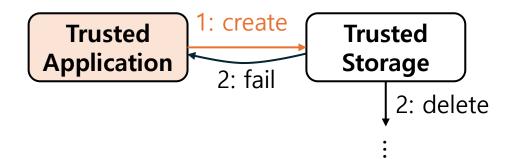


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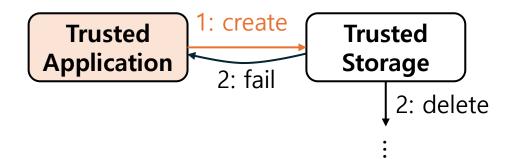


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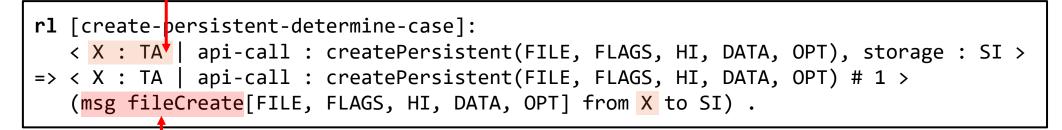
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TA

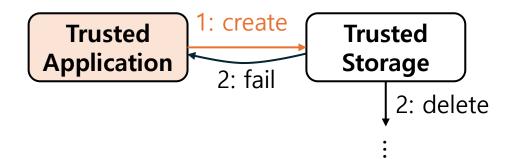


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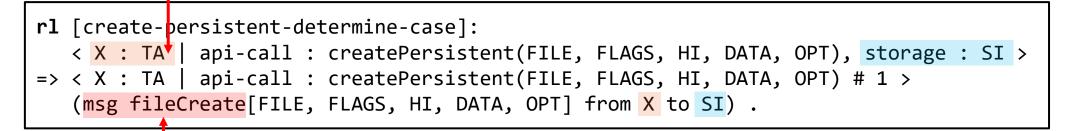


Make a file creation request message and send it

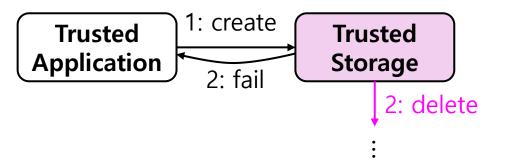


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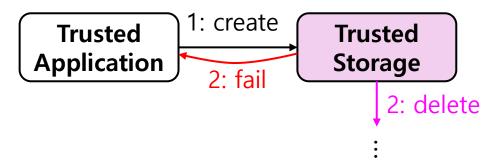
TA



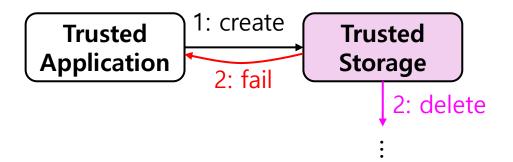
Make a file creation request message and send it to its trusted storage



• (2)-1. The storage deletes the old file if an overwrite flag is given.

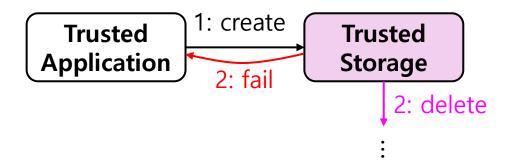


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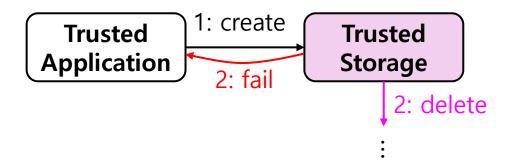


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- Trusted storage has the following things:
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- ...

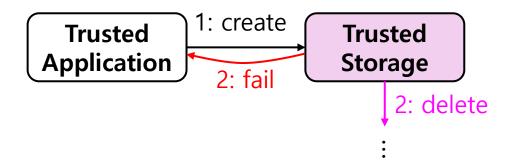
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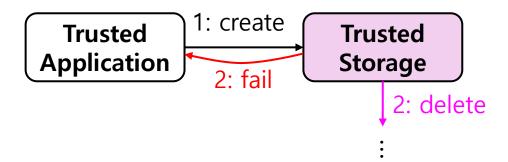


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```
- a || class Storage | files : Set{FileName}, counter : Nat, ...
```

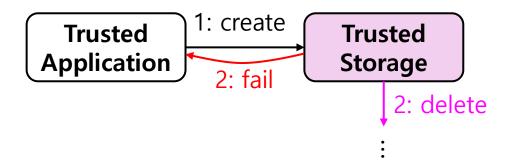
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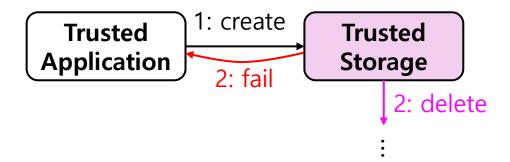


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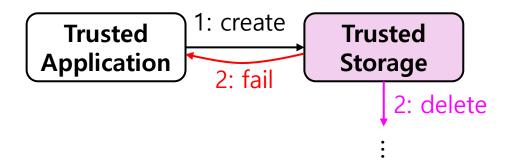
```
crl [create-persistent-overwrite-check]:
   (msg create[METHOD FILE FLAGS HI DATA] from X to SI)
   < PI : PersistObj | file-name : FILE >
        < SI : Storage | status : normal, files : FILES, counter : N >
=> < PI : PersistObj | >
        if overwrite in FLAGS
        then < SI : Storage | counter : N + 2 >
            (msg create[METHOD FILE FLAGS HI DATA N X] from SI to PI)
        else (msg createFail from SI to TK) < SI : Storage | > fi if FILE in FILES .
```



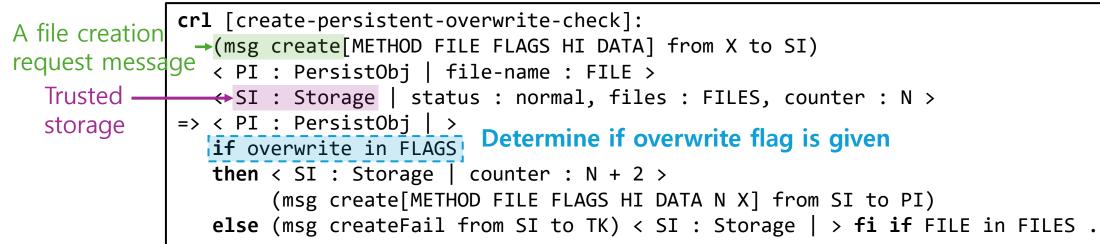
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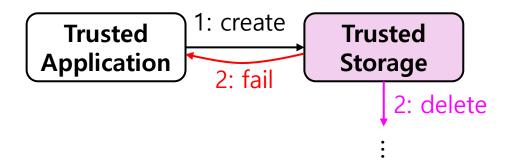


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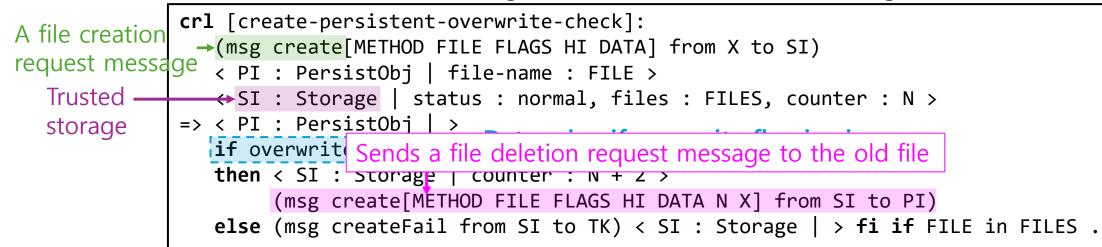


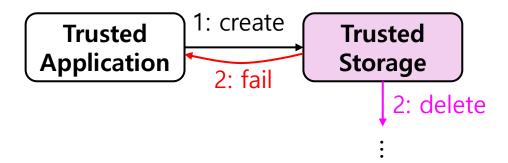
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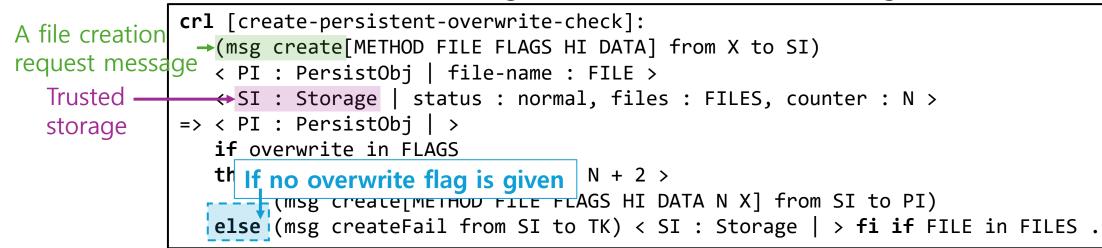


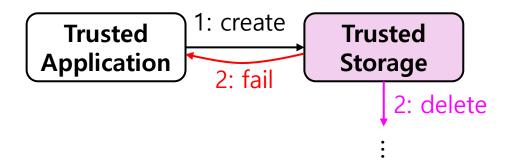
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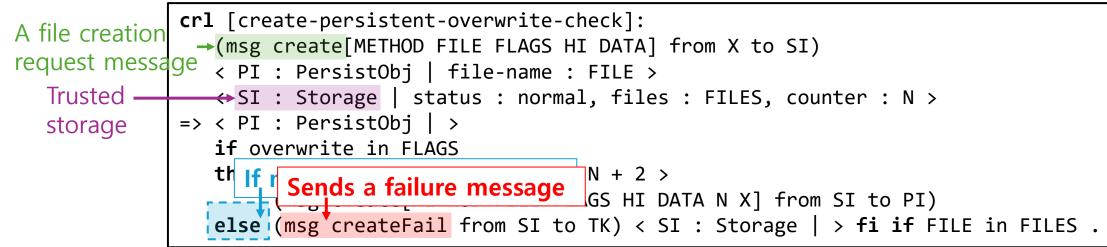


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Trusted Storage API (27/27)

TEE_CreatePersistentObject TEE_OpenPersistentObject TEE_RenamePersistentObject TEE_CloseAndDeletePersistentObject1 TEE_ReadObjectData TEE_WriteObjectData

... TEE_CopyObjectAttributes1 TEE_PopulateTransientObject

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... TEE_CopyObjectAttributes1 TEE_PopulateTransientObject Crytographic Operations API (30/30)

TEE_AllocateOperation TEE_ResetOperation TEE_SetOperationKey TEE_CopyOperation TEE_FreeOperation TEE_DigestUpdate

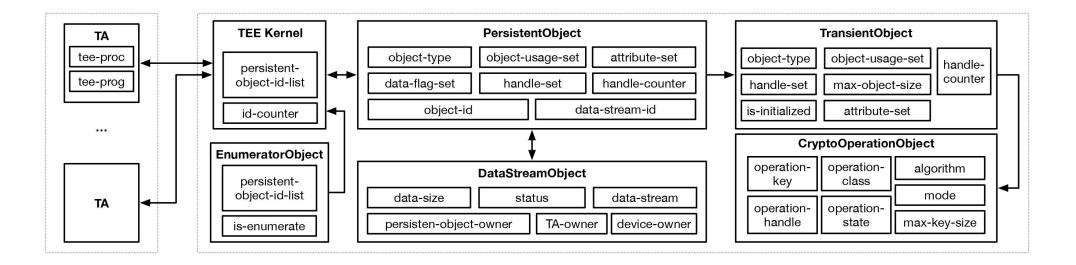
... TEE_MACInit TEE_MACUpdate

...

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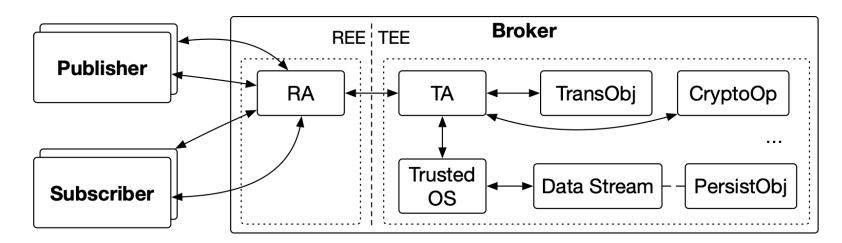
Settings

- We define the language semantics for TEE applications in Maude.
- We extend our model to run TEE applications using this semantics.

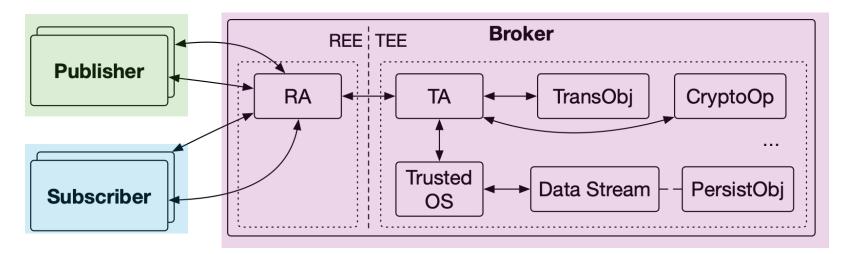
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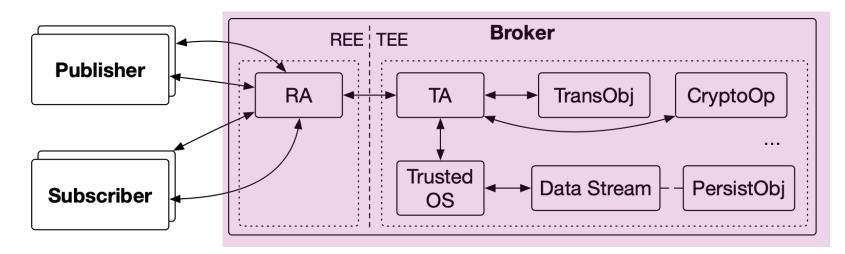
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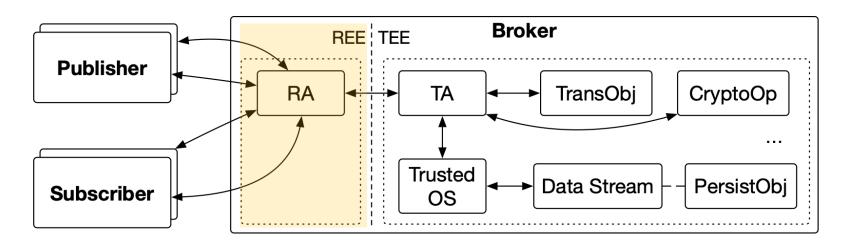
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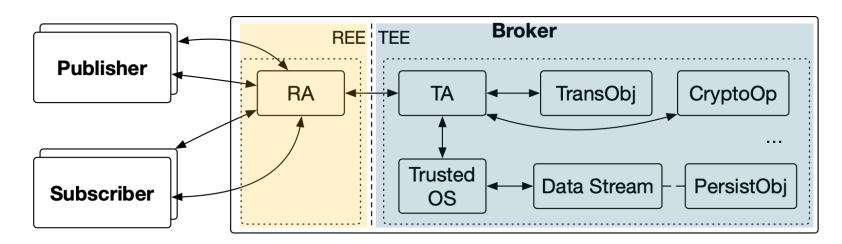
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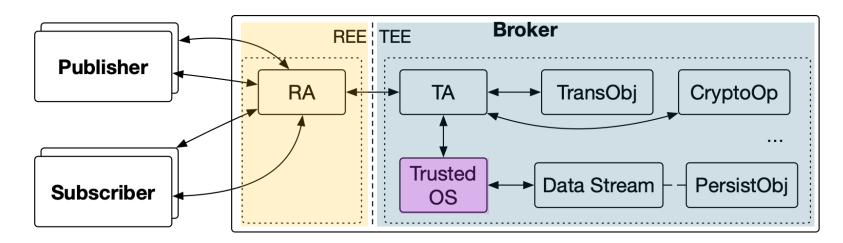
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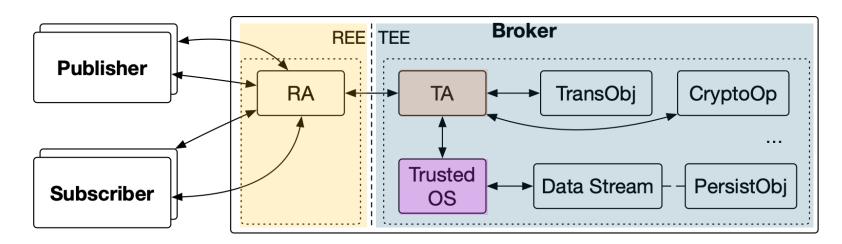
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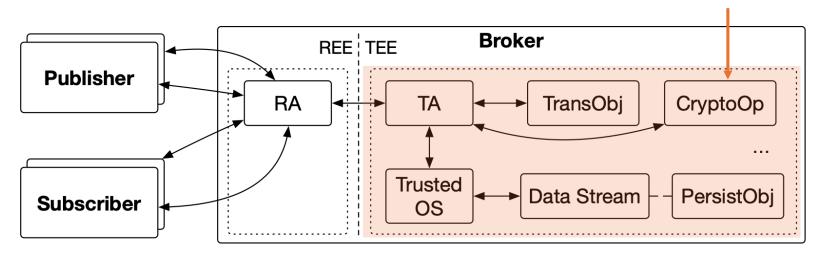
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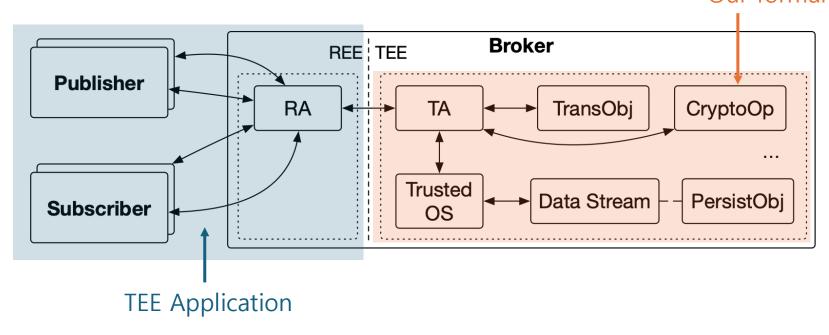
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Threat Models

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 - This threat makes brokers to run out of memory.

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- (2) Message modification threat
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Name	Description	LTL Formula
P1	If no memory error occurs in the broker, subscribers eventually receive messages.	$\Box \neg memErr.B \rightarrow \\ \Box (send.P \rightarrow \diamond recv.S)$
P2	If the TA panics, subscribers should not receive any messages.	$\Box (panic.TA \rightarrow \Box \neg recv.S)$
P3	If any memory error occurs in the broker, subscribers should not receive any messages.	$\Box (memErr.B)$ $\rightarrow \Box \neg recv.S)$
P4	When the TA starts running, it should eventually terminate.	$\Box (start.TA \rightarrow term.TA)$
P5	If subscribers receive messages from publishers, messages sent from each publisher are in order.	$\Box (inQueue. P(a :: b :: c)) \rightarrow \diamond inQueue. S(a :: b :: c)$
P6	The number of tasks handled by the TA cannot exceed five.	\Box (¬ <i>numTaskExceed</i> (5))

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	NON	Т	62	35.7		NON	Т	62	35		NON	Т	62	33.8
P1	MSG	Т	148	90.1	P3	MSG	Т	148	88.8	P5	MSG	Т	148	86.9
	OOM	Т	202	144.2		OOM	\perp	0.1	0.1		OOM	Т	532	546.7
	NON	T	62	34.9		NON	Т	62	34.9		NON	Т	62	34.3
P2	MSG	\perp	17	9.1	P4	MSG	Т	148	88.6	P6	MSG	Т	148	87.9
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P3	If any receive		-			irs in t	the bro	oker,	subsc	ribe	rs sho	uld no	ot			•	nErr.B :ecv.S)	
	Prop. Type Safe? S Time Prop. Type Safe? S Time Prop. Type Sa																	
]	Prop.	Type	Safe?	$ \mathbf{S} $	Time	Prop.	Type	Safe?	$ \mathbf{S} $	Time	Prop.	Type	Safe?	$ \mathbf{S} $	Time		
			NON	Т	62	35.7		NON	Т	62	35		NON	Т	62	33.8		
		P1	MSG	Т	148	90.1	P3	MSG		148	88.8	P5	MSG	Т	148	86.9		
			OOM	Т	202	144.2		OOM	\perp	0.1	0.1		OOM	Т	532	546.7		
			NON	Т	62	34.9		NON	T	62	34.9	•	NON	Т	62	34.3		
		P2	MSG		17	9.1	P4	MSG	Т	148	88.6	P6	MSG	Т	148	87.9		
	_		OOM	Τ	532	547.9		OOM	Т	532	539.3		OOM	Т	532	542.4		

P2 If Even if the	e TA pan	icked, so	me suł	oscribe	r recei	ves a	messag	e. 🗆	(pani	c.TA –	$\rightarrow \Box \neg recv.S)$
P3 If If memory receive any r			n TA, s	ome si	ubscrib	er sti	ll receive	es a m		ge. m	nErr.B recv.S)
Prop. T	ype Safe?	$ \mathbf{S} $ Time	Prop. 7	Type Sa	fe? S	Time	Prop. Ty	rpe Safe	$ \mathbf{S} $	Time	
	NON T	62 35.7		NON	\top 62	35		DN ⊤	62		
	1SG ⊤	148 90.1	P3	MSG	T 148	88.8	1	$G \top$	148		
C	DOM	$202 \ 144.2$		OOM _	$\perp 0.1$	0.1	00	DM	532	546.7	
٨	NON T	62 34.9		NON	\top 62	34.9	N	ON ⊤	62	34.3	
P2 N	1SG ⊥	17 9.1	P4	MSG	\top 148	88.6	P6 M3	SG ⊤	148	87.9	
(ООМ Т	$532 \ 547.9$	·	OOM	\top 532	539.3	0	OM ⊤	532	542.4	

- The reason is that the broker program cannot distinguish the following three TA status:
 - (1) successful termination,
 - (2) panic,
 - (3) out-of-memory.

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• We propose a code-level patch for the broker program to distinguish two error states from successful termination.

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```
{ . . .
```

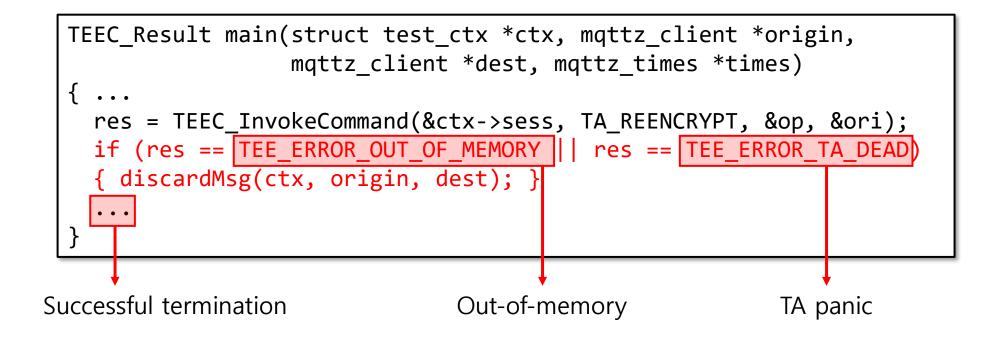
. . .

res = TEEC_InvokeCommand(&ctx->sess, TA_REENCRYPT, &op, &ori);

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• We propose a code-level patch for the broker program to distinguish two error states from successful termination.

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	NON	Т	62	35.3		NON	Т	62	34.8		NON	Т	62	34.1
P1	MSG	Т	149	89.9	P3	MSG	Т	149	89.7	P5	MSG	\top	149	87.4
	OOM	Т	203	146.2		OOM	Т	347	285.2		OOM	Т	347	288.6
	NON	Т	62	35.1		NON	Т	62	34.7		NON	T	62	34.4
P2	MSG	Т	149	89.9	P4	MSG	Т	149	89.4	P6	MSG	Т	149	87.9
	OOM	Т	347	294.8		OOM	Т	347	278.5		OOM	Т	347	286.1

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P1	MSG	Т	149	89.9	P3	MSG	Т	149	89.7	P5	MSG	Т	149	87.4
	OOM	Т	203	146.2		OOM	Т	347	285.2		OOM	Т	347	288.6
	NON	Т	62	35.1		NON	Т	62	34.7		NON	Т	62	34.4
P2	MSG	Т	149	89.9	P4	MSG	Т	149	89.4	P6	MSG	Т	149	87.9
	OOM	Т	347	294.8		OOM	Т	347	278.5		OOM	Т	347	286.1

We can confirm that the violated properties are satisfied.

Summary

- We provide a comprehensive formal model for <u>TEE APIs</u>, that can be used in various formal analysis.
- We specify two widely used TEE API categories, <u>Trusted Storage API</u> and <u>Cryptographic Operations API</u>.
- We demonstrate the effectiveness of our model through a case study on formally analyzing <u>a real-world TEE application</u>, MQT-TZ.